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Reg. No.....

Name.....

**M.Sc. DEGREE (C.S.S.) EXAMINATION, MAY 2020**

**Fourth Semester**

Faculty of Science

Branch I (A) : Mathematics

MT 04 E14—CODING THEORY

(2012 Admission onwards)

Time : Three Hours

Maximum Weight : 30

**Part A**

*Answer any five questions.*

*Each question has weight 1.*

1. Show that the set of all binary words of length  $n$  is a vector space.
2. Compute the weight of each of the following words and the distance between each pair of them.  
 $V_1 = 1001010$ ,  $V_2 = 0110101$ ,  $V_3 = 0011110$ .
3. Show that if  $C$  is a binary  $[n, (n - 1)/2]$  self orthogonal code, for odd  $n$ , then  $C^+$  is the  $[n, (n + 1)/2]$  code generated by  $C$  and  $h$ .
4. Show that a self dual  $\left[ n, \frac{n}{2} \right]$  ternary code exists iff  $n$  is divisible by 4.
5. Show that the order of any element  $g \in G$ , divides the order of  $G$ , where  $G$  is a finite group.
6. Find the minimal polynomials for the elements 0 and 1 in  $GF(16)$  constructed using  $1 + x + x^3$ .
7. Show that if a binary cyclic code with generator polynomial  $g(x)$  is self-orthogonal then  $1 + x$  must divide  $g(x)$ .
8. Prove that a Reed-Solomon code  $C$  of designed distance  $d$  has  $d$  as its actual minimum weight. Also show that  $C$  is an MDS code.

(5 × 1 = 5)

**Turn over**





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### Part B

*Answer any five questions.  
Each question has weight 2.*

9. Prove that if the rows of a generator matrix  $G$  for a binary  $[n, k]$  code  $C$  have weights divisible by 4 and are orthogonal to each other, then  $C$  is self-orthogonal and all weights in  $C$  are divisible by 4.
10. Define syndrome. Show that if  $C$  is a binary code and  $e$  is any vector, the syndrome of  $e$  is the sum of those columns of  $H$  where  $e$  has non-zero components ; where  $H$  be a parity check matrix of an  $[n, k]$  code.
11. Define Golay code. Find the minimum weight and show that it is triple error correcting code.
12. Show that every monic polynomial over a field  $F$  can be expressed uniquely as a product of irreducible monic polynomials over  $F$ .
13. Let  $x^n - 1 = g(x)h(x)$  over  $GF(q)$ . Prove that a cyclic code  $C$  with generator polynomial  $g(x)$  is self-orthogonal iff the reciprocal polynomial of  $h(x)$  divides  $g(x)$ .
14. Find a self-orthogonal length 15 binary cyclic code.
15. Show that in a field of characteristics  $P(x \pm y)^{P^m} = x^{P^m} \pm y^{P^m}$ .
16. Show that for any prime  $p$  and positive integer  $m$ , there is a unique field of  $p^m$  elements.

(5 × 2 = 10)

### Part C

*Answer any three questions.  
Each question has weight 5.*

17. (a) If  $u$  is a vector in  $c$  of weight  $s$ , show that there is a dependence relation among  $s$  columns of any parity check matrix of  $c$  and conversely that any dependence relation among  $s$  columns of a parity check matrix of  $c$  yields a vector of weight  $s$  in  $c$ .
- (b) If  $C$  has minimum weight  $d$ , show that  $c$  can detect all errors of weight  $\leq d - 1$ .
- (c) If  $C$  has minimum weight  $2(t + 1)$ , show that we can simultaneously correct all errors of weight  $t$  or less and detect all errors of weight  $t + 1$ .





18. Prove the following :

- (a) (i) Given  $m$  and  $d$ , then there exists a binary code of length  $n$ , minimum distance  $d$  or more and dimension  $k \geq n - m$ , whenever

$$\binom{n-1}{1} + \binom{n-1}{2} + \dots + \binom{n-1}{d-2} < 2^m - 1.$$

- (ii) Given  $m$  and  $d$ , then there exists a code over  $\text{GF}(q)$  of length  $n$ , minimum distance  $d$  or more and dimension  $k \geq n - m$ , whenever

$$(q-1)\binom{n-1}{1} + (q-1)^2\binom{n-1}{2} + \dots + (q-1)^{d-2}\binom{n-1}{d-2} < q^m - 1.$$

- (b) If  $d$  is even show that  $A(n-1, d-1) = A(n, d)$ .

19. (a) Messages are encoded using  $C_{15}$ . Determine if possible the location of the errors if  $w$  is received with syndrome  $wH$  is given by 01000000. The parity check matrix  $H$  for  $C_{15}$  is given below :

$$H = \begin{bmatrix} 1000 & 1000 \\ 0100 & 0001 \\ 0010 & 0011 \\ 0001 & 0101 \\ 1100 & 1111 \\ 0110 & 1000 \\ 0011 & 0001 \\ 1101 & 0011 \\ 1010 & 0101 \\ 0101 & 1111 \\ 1110 & 1000 \\ 0111 & 0001 \\ 1111 & 0011 \\ 1011 & 0101 \\ 1001 & 1111 \end{bmatrix}$$

- (b) Prove that a doubly even  $\left[ n, \frac{n}{2} \right]$  code exists iff  $n$  is divisible by 8.

Turn over





20. (a) Show that  $F(x)/(f(x))$  is a field iff  $f(x)$  is irreducible. Also show that if  $f(x)$  be an irreducible polynomial of degree  $m$  over  $GF(p)$ , then  $F' = GF(p)[x]/(f(x))$  is a field with  $p^m$  elements.
- (b) Prove that every finite field has a primitive element.
21. (a) If  $g(x)h(x) = x^n - 1$  in  $F[x]$  and  $g(x)$  is the generator polynomial of a cyclic code  $C$ , show that the reciprocal polynomial of  $h(x)$  is the generator polynomial of  $C^\perp$ . Also show that if  $h(x) = h_0 + h_1(x) + \dots + h_k(x^k)$ , then the matrix  $H$  given below is a parity check matrix of  $C$ , and  $C^\perp$  is cyclic.

$$H = \begin{pmatrix} h_k & h_{k-1} \dots & h_0 & 00 \dots 0 \\ 0 & h_k \dots & h_1 & h_0 0 \dots 0 \\ " & " & " & \\ " & " & " & \\ 0 & 0 \dots & h_k \dots & h_0 \end{pmatrix}$$

- (b) Prove that every cyclic  $[n, k]$  code  $C$  has an idempotent generator  $e(x)$ .
22. (a) Prove that the minimum weight of a BCH code  $C$  of designed distance  $\sigma$  is at least  $\sigma$ .
- (b) Let  $f(x)$  be a polynomial with co-efficients in  $GF(q)$  and let  $S$  be the set of its roots in some field  $F = GF(q^m)$ . Prove that the weight of is greater then or equal to the size of any set  $A$  in a set  $I_s$  of subsets of  $F$  that is independent with respect to  $S$ .

(3 × 5 = 15)

